

Derivatives of WS1S Formulas

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Logic-Automaton Connection

WS1S

$\varphi = T \mid F \mid x \in X \mid x < y \mid \varphi \vee \psi \mid \neg \varphi \mid \exists x. \varphi \mid \exists X. \varphi$
finite

Logic-Automaton Connection

WS1S

$\forall I. I \models \varphi \iff I \models \psi?$

$\varphi = T \mid F \mid x \in X \mid x < y \mid \varphi \vee \psi \mid \neg \varphi \mid \exists x. \varphi \mid \exists X. \varphi$
finite

Logic-Automaton Connection

Finite Automata

↑
MONA

WS1S

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Logic-Automaton Connection

Finite Automata

MONA

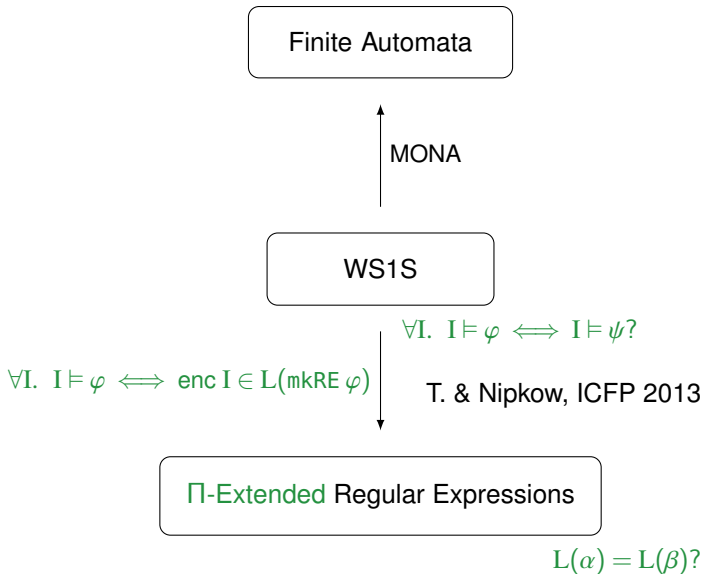
WS1S

$\forall I. I \models \varphi \iff I \models \psi?$

Regular Expressions

$L(\alpha) = L(\beta)?$

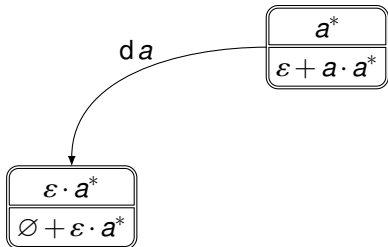
Logic-Automaton Connection



$$a^* \stackrel{?}{\equiv} \varepsilon + a \cdot a^* \text{ for } \Sigma = \{a, b\}$$

a^*
$\varepsilon + a \cdot a^*$

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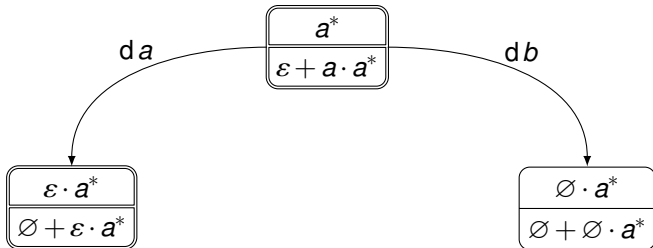


Brzozowski derivative

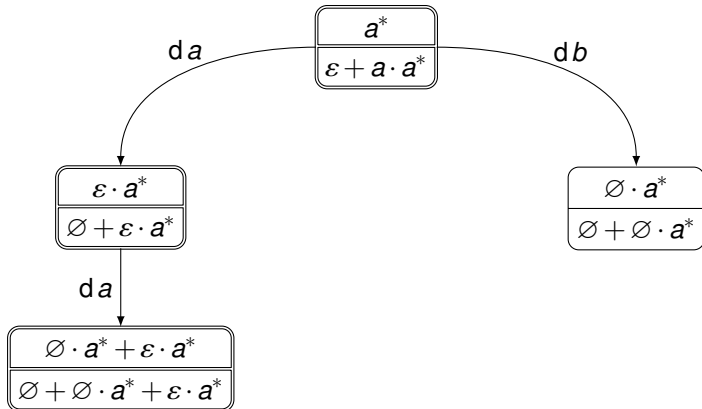
d : letter \rightarrow regex \rightarrow regex

$$L(dar) = \{w \mid aw \in L(r)\}$$

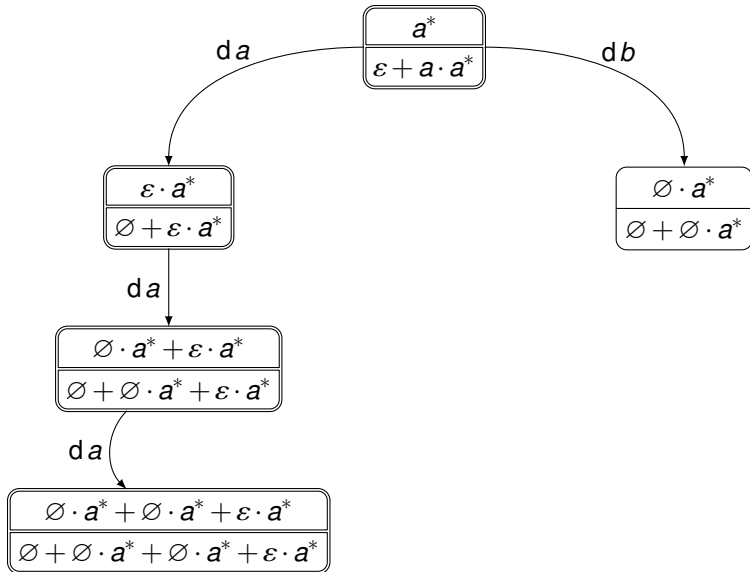
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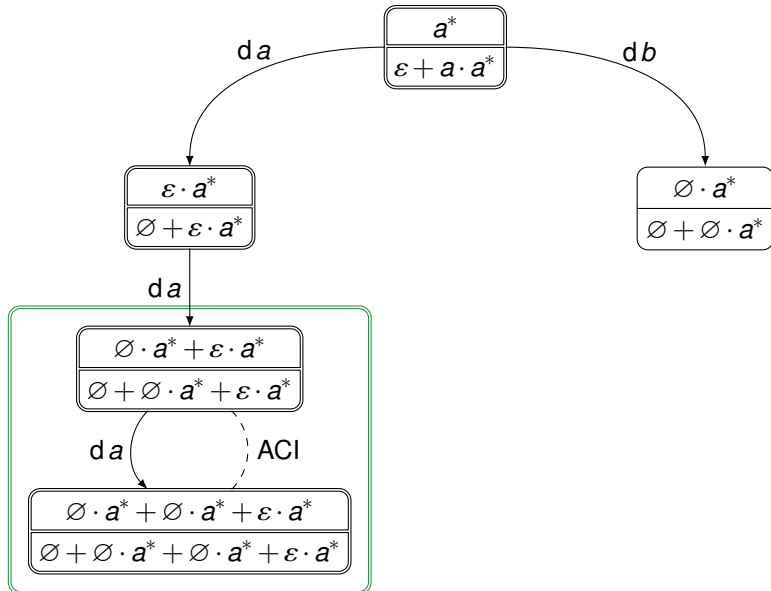
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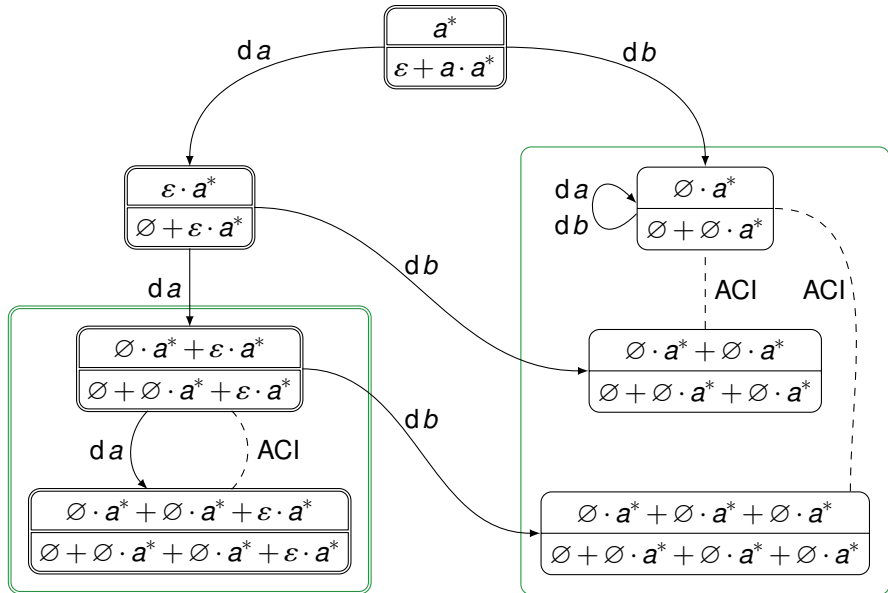
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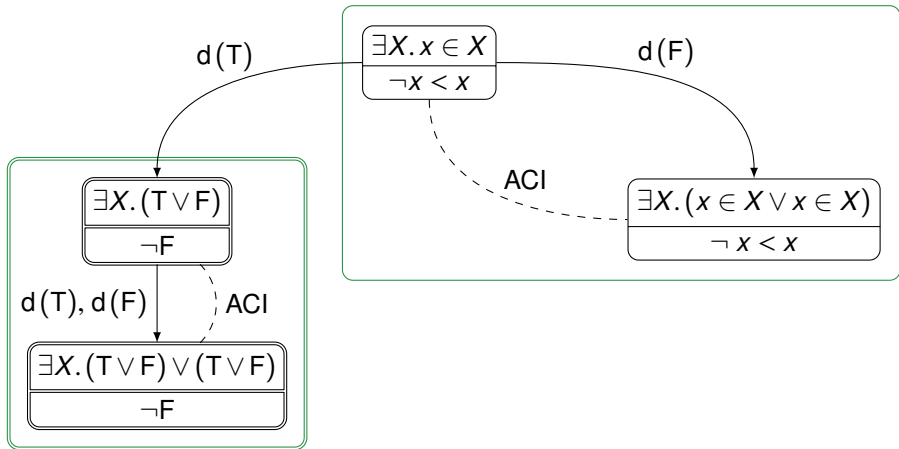


Key ingredients: Derivative + Acceptance Test

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Let's define them on WS1S formulas directly!

$(\exists X. x \in X) \stackrel{?}{\equiv} (\neg x < x)$ for $\Sigma = \{(T), (F)\}$



Derivative

$$d v T = T$$

$$d v F = F$$

$$d v (x \in X) = \begin{cases} x \in X & \text{if } \neg v[x] \\ T & \text{if } v[x] \wedge v[X] \\ F & \text{otherwise} \end{cases}$$

$$d v (x < y) = \dots$$

$$d v (\varphi \vee \psi) = d v \varphi \vee d v \psi$$

$$d v (\neg \varphi) = \neg d v \varphi$$

$$d v (\exists X. \varphi) = \exists X. (d (v_{X \mapsto T}) \varphi \vee d (v_{X \mapsto F}) \varphi)$$

Acceptance Test

$$\begin{aligned}\varepsilon T &= T \\ \varepsilon F &= F \\ \varepsilon (x \in X) &= F \\ \varepsilon (x < y) &= F \\ \varepsilon (\varphi \vee \psi) &= \varepsilon \varphi \vee \varepsilon \psi \\ \varepsilon (\neg \varphi) &= \neg \varepsilon \varphi \\ \varepsilon (\exists X. \varphi) &= \text{action happens here}\end{aligned}$$

Acceptance Test

εT	=	T
εF	=	F
$\varepsilon (x \in X)$	=	F
$\varepsilon (x < y)$	=	F
$\varepsilon (\varphi \vee \psi)$	=	$\varepsilon \varphi \vee \varepsilon \psi$
$\varepsilon (\neg \varphi)$	=	$\neg \varepsilon \varphi$
$\varepsilon (\exists X. \varphi)$	=	action happens here

futurization

derivatives
from the right

Altogether

A decision procedure for **WS1S** that

Altogether

A decision procedure for **WS1S** that
operates **on formulas** directly and

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A decision procedure for **WS1S** that
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Thanks. Questions?

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