

# Semantics of Programming Languages

## Exercise Sheet 1

Before beginning to solve the exercises, open a new theory file named `Ex01.thy` and write the the following three lines at the top of this file.

```
theory Ex01
imports Main
begin
```

### Exercise 1.1 Calculating with natural numbers

Use the `value` command to turn Isabelle into a fancy calculator and evaluate the following natural number expressions:

`"2 + (2::nat)"`      `"(2::nat) * (5 + 3)"`      `"(3::nat) * 4 - 2 * (7 + 1)"`

Can you explain the last result?

### Exercise 1.2 Natural number laws

Formulate and prove the well-known laws of commutativity and associativity for addition of natural numbers.

### Exercise 1.3 Counting elements of a list

Define a function which counts the number of occurrences of a particular element in a list.

```
fun count :: "'a list ⇒ 'a ⇒ nat"
```

Test your definition of `count` on some examples and prove that the results are indeed correct.

Prove the following inequality (and additional lemmas, if necessary) about the relation between `count` and `length`, the function returning the length of a list.

```
theorem "count xs x ≤ length xs"
```

### Exercise 1.4 Adding elements to the end of a list

Recall the definition of lists from the lecture. Define a function *snoc* that appends an element at the right end of a list. Do not use the existing append operator `@` for lists.

```
fun snoc :: "'a list  $\Rightarrow$  'a  $\Rightarrow$  'a list"
```

Convince yourself on some test cases that your definition of *snoc* behaves as expected, for example run:

```
value "snoc [] c"
```

Also prove that your test cases are indeed correct, for instance show:

```
lemma "snoc [] c = [c]"
```

Next define a function *reverse* that reverses the order of elements in a list. (Do not use the existing function *rev* from the library.) Hint: Define the reverse of  $x \# xs$  using the *snoc* function.

```
fun reverse :: "'a list  $\Rightarrow$  'a list"
```

Demonstrate that your definition is correct by running some test cases, and proving that those test cases are correct. For example:

```
value "reverse [a, b, c]"
```

```
lemma "reverse [a, b, c] = [c, b, a]"
```

Prove the following theorem. Hint: You need to find an additional lemma relating *reverse* and *snoc* to prove it.

```
theorem "reverse (reverse xs) = xs"
```

### Homework 1 Sum of odd numbers

*Submission until Wednesday, November 2, 12:00 (noon).*

In this homework assignment you will prove that the square of a natural number  $n$  can be computed as the sum of the first  $n$  odd numbers, which we will write as *oddsun*( $n$ ). For example, we have  $\text{oddsun}(3) = 1 + 3 + 5 = 9 = 3 * 3$ .

Your first task is to use the **fun** command to define a function  $\text{oddsun} :: \text{nat} \Rightarrow \text{nat}$  in Isabelle. Your definition should have equations for  $\text{oddsun}(0)$  and  $\text{oddsun}(\text{Suc } n)$ .

```
fun oddsun :: "nat  $\Rightarrow$  nat"
```

You may wish to use the **value** command to check that your definition is correct. For example, the following command should evaluate to *True*.

```
value "oddsun 3 = 1 + 3 + 5"
```

Your next task is to prove by induction that for any  $n$ , *oddsun*( $n$ ) computes the square of  $n$ . First, write an informal proof by hand. Your proof should contain a base case for

zero, where you show that  $oddsun(0)$  equals the square of  $0$ . Next you should have a case for successor: Fix an arbitrary  $m$ , assume the inductive hypothesis that  $oddsun(m)$  equals the square of  $m$ , and then show that  $oddsun(Suc\ m)$  equals the square of  $Suc\ m$ .

Finally, prove the same property *formally* in Isabelle:

**lemma** " $oddsun\ n = n * n$ "