

# Semantics of Programming Languages

## Exercise Sheet 8

### Exercise 8.1 Definite Assignment Analysis

In the lecture, you have seen a definite assignment analysis that was based on the large-step semantics. Definite assignment analysis can also be based on a small-step semantics. Furthermore, the ternary predicate  $D$  from the lecture can be split into two parts: a function  $AA :: com \Rightarrow name\ set$  (“assigned after”) which collects the names of all variables assigned by a command and a binary predicate  $D :: name\ set \Rightarrow com \Rightarrow bool$  which checks that a command accesses only previously assigned variables. Conceptually, the ternary predicate from the lecture (call it  $D_{lec}$ ) and the two-step approach should relate by the equivalence  $D\ V\ c \longleftrightarrow D_{lec}\ V\ c\ (V \cup AA\ c)$

1. Download the theory `ex08.template` and study the already defined small-step semantics for definite analysis.
2. Define the function  $AA$  which computes the set of variables assigned after execution of a command. Furthermore, define the predicate  $D$  which checks if a command accesses only assigned variables, assuming the variables in the argument set are already assigned.
3. Prove progress and preservation of  $D$  with respect to the small-step semantics, and conclude soundness of  $D$ . You may use (and then need to prove) the lemmas  $D\_incr$  and  $D\_mono$ .

### Homework 8 Read Variables

*Submission until Wednesday, December 21, 2011, 12:00 (noon).*

Instantiates the `vars` typeclass for commands, such that  $vars\ c$  is the set of variables read by the command.

Then show, that an execution does not depend on variables not read by the command, w.r.t. the small-step semantics. I.e., show the following lemma:

**lemma** “ $\llbracket (c,s) \rightarrow^* (c',s'); s = t\ on\ X; vars\ c \subseteq X \rrbracket$   
 $\implies \exists t'. (c,t) \rightarrow^* (c',t') \wedge s' = t'\ on\ X$ ”

Hint: You may want to show the lemma for a single small-step first, i.e.,

**lemma** *eq\_step*: “ $\llbracket (c,s) \rightarrow (c',s'); s = t \text{ on } X; \text{ vars } c \subseteq X \rrbracket$   
 $\implies \exists t'. (c,t) \rightarrow (c',t') \wedge s' = t' \text{ on } X$ ”