# Semantics of Programming Languages

Exercise Sheet 13

### Exercise 13.1 Sign Analysis

Instantiate the abstract interpretation framework to a sign analysis over the lattice *pos, zero, neg, any*, where *pos* abstracts positive values, *zero* abstracts zero, *neg* abstracts negative values, and any abstracts any value.

For this exercise, you best modify the parity analysis src/HOL/IMP/Abs\_Int1\_parity.

### Exercise 13.2 Complete Lattices: GLB of UBs is LUB

Show that the least upper bound is the greatest lower bound of all upper bounds: definition "Sup'  $(S::'a::complete\_lattice set) \equiv Inf \{u. \forall s \in S. s \leq u\}$ "

lemma  $Sup'\_upper$ : " $\forall s \in S. s \leq Sup' S$ " lemma  $Sup'\_least$ : assumes upper: " $(\forall s \in S. s \leq u)$ " shows " $Sup' S \leq u$ "

## Homework 13.1 Lattice Theory

Submission until Tuesday, 07.02.2017 (Pen & Paper), 10:00am.

A type 'a is a  $\sqcup$ -semilattice if it is a partial order and there is a supremum operation  $\sqcup$  of type 'a  $\Rightarrow$  'a  $\Rightarrow$  'a that returns the least upper bound of its arguments:

- Upper bound:  $x \leq x \sqcup y$  and  $y \leq x \sqcup y$
- Least:  $x \leq z \land y \leq z \longrightarrow x \sqcup y \leq z$

Is every finite  $\sqcup$ -semilattice with a bottom element  $\bot$  also a complete lattice? Prove or disprove on paper!

#### Homework 13.2 Al for the Extended Reals

Submission until Tuesday, 07.02.2017 (Isabelle), 10:00am. For this exercise, we will consider a modified variant of IMP that computes on real numbers extended with  $-\infty$  and  $\infty$ . The corresponding type is *ereal*. We will consider " $-\infty + \infty$ " and " $\infty + (-\infty)$ " erroneous computations. We propagate errors by using the option type, i.e. we set  $val = ereal \ option$ . The files AExp.thy, BExp.thy, BigStep.thy, Collecting.thy for this variant are provided for you on the website. Your task is now to design an abstract interpreter on the domain consisting of subsets of  $\{\infty^-, \infty^+, NaN, Real\}$  where NaN signals a computation error and all other values have their obvious meaning. First adopt  $Abs\_Int0.thy$  and  $Abs\_Int1.thy$  to accommodate for the changed semantics, and then instantiate the abstract interpreter ( $Abs\_Int, Abs\_Int\_mono$ , and  $Abs\_Int\_parity.thy$ . Hints: To benefit from proof automation it can be helpful to slightly change the format of the rules for addition in  $Val\_semilattice$ . For instance, you could reformulate  $gamma\_plus'$  as:  $i1 \in \gamma \ a1 \implies i2 \in \gamma \ a2 \implies i = i1 + i2 \implies i \in \gamma(plus' \ a1 \ a2)$ . (You will need to change the interface  $Val\_semilattice$ ).

You can start the formalization of the AI like this:

datatype bound = NegInf (" $\infty^-$ ") | PosInf (" $\infty^+$ ") | NaN | Real

**datatype** bounds = S "bound set" instantiation bounds :: order begin definition less\_eq\_bounds where " $x \le y = (case \ (x, y) \ of \ (S \ x, S \ y) \Rightarrow x \subseteq y)$ " definition less\_bounds where " $x < y = (case \ (x, y) \ of \ (S \ x, S \ y) \Rightarrow x \subset y)$ "