

# Semantics of Programming Languages

## Exercise Sheet 12

### Exercise 12.1 Complete Lattices

Which of the following ordered sets are complete lattices?

- $\mathbb{N}$ , the set of natural numbers  $\{0, 1, 2, 3, \dots\}$  with the usual order.
- A finite set  $A$  with a total order  $\leq$  on it.
- $\mathbb{N} \cup \{\infty\}$ , the set of natural numbers plus infinity, with the usual order and  $n < \infty$  for all  $n \in \mathbb{N}$ .

### Exercise 12.2 Sign Analysis

Instantiate the abstract interpretation framework to a sign analysis over the lattice  $pos, zero, neg, any$ , where  $pos$  abstracts positive values,  $zero$  abstracts zero,  $neg$  abstracts negative values, and  $any$  abstracts any value.

```
datatype sign = Pos | Zero | Neg | Any

instantiation sign :: order
instantiation sign :: semilattice_sup_top
fun  $\gamma\_sign$  :: "sign  $\Rightarrow$  val set"
fun num_sign :: "val  $\Rightarrow$  sign"
fun plus_sign :: "sign  $\Rightarrow$  sign  $\Rightarrow$  sign"
global\_interpretation Val_semilattice
  where  $\gamma$  =  $\gamma\_sign$  and num' = num_sign and plus' = plus_sign
global\_interpretation Abs_Int
  where  $\gamma$  =  $\gamma\_sign$  and num' = num_sign and plus' = plus_sign
  defines aval_sign = aval' and step_sign = step' and AI_sign = AI
```

Some tests:

```
definition "test1_sign =
  'x' ::= N 1;;
  WHILE Less (V 'x') (N 100) DO 'x' ::= Plus (V 'x') (N 2)"
value "show_acom (the(AI_sign test1_sign))"
```

```
definition "test2_sign =
  'x' ::= N 1;;
```

*WHILE Less (V "x") (N 100) DO "x" ::= Plus (V "x") (N 3)*

**definition** “*steps c i = ((step\_sign  $\top$ )  $\sim$  i) (bot c)*”

**value** “*show\_acom (steps test2\_sign 0)*”

...

**value** “*show\_acom (steps test2\_sign 6)*”

**value** “*show\_acom (the(AI\_sign test2\_sign))*”

### Exercise 12.3 AI for Conditionals

Our current constant analysis does not regard conditionals. For example, it cannot figure out, that after executing the program *x:=2; IF x<2 THEN x:=2 ELSE x:=1*, *x* will be constant.

In this exercise, we extend our abstract interpreter with a simple analysis of boolean expressions. To this end, modify locale *Val\_semilattice* from theory *Abs\_Int0.thy* as follows:

- Introduce an abstract domain *'bv* for boolean values, add, analogously to *num'* and *plus'* also functions for the boolean operations and for *less*.
- Modify *Abs\_Int0* to accommodate for your changes.

### Homework 12.1 AI for the Extended Reals

*Submission until Wednesday, Jan 22, 23:59pm.* For this exercise, we will consider a modified variant of IMP that computes on real numbers extended with  $-\infty$  and  $\infty$ . The corresponding type is *ereal*. We will consider “ $-\infty + \infty$ ” and “ $\infty + (-\infty)$ ” erroneous computations. We propagate errors by using the *option* type, i.e. we set *val = ereal option*. The theories up to *Collecting* for this variant are already provided. Your task is now to design an abstract interpreter on the domain consisting of subsets of  $\{\infty^-, \infty^+, NaN, Real\}$  where *NaN* signals a computation error and all other values have their obvious meaning. First adopt *Abs\_Int0* and *Abs\_Int1* to accommodate for the changed semantics, and then instantiate the abstract interpreter with your analysis. For this step you best modify the parity analysis *Abs\_Int1\_parity*.

*Hints:* To benefit from proof automation it can be helpful to slightly change the format of the rules for addition in *Val\_semilattice*. For instance, you could reformulate *gamma\_plus'* as:  $i1 \in \gamma a1 \implies i2 \in \gamma a2 \implies i = i1 + i2 \implies i \in \gamma(plus' a1 a2)$ . (You will need to change the interface *Val\_semilattice*).

You can start the formalization of the AI like this:

**datatype** *bound* =  $\infty^-$  |  $\infty^+$  | *NaN* | *Real*

**datatype** *bounds* = *B* (*bound set*)

**instantiation** *bounds* :: *order*

**begin**

**definition** *less\_eq\_bounds* **where**

“ $x \leq y = (\text{case } (x, y) \text{ of } (B\ x, B\ y) \Rightarrow x \subseteq y)$ ”

**definition** *less\_bounds* **where**

“ $x < y = (\text{case } (x, y) \text{ of } (B\ x, B\ y) \Rightarrow x \subset y)$ ”

**instance**

**end**

For the AI, interpret *Abs\_Int*, *Abs\_Int\_mono*, and *Abs\_Int\_measure*:

**instantiation** *bounds* :: *semilattice\_sup\_top*

**begin**

**definition** *sup\_bounds*

**definition** *top\_bounds*

**instance**

**end**

**fun** *γ\_bounds* :: “*bounds*  $\Rightarrow$  *val set*”

**definition** *num\_bounds* :: “*ereal*  $\Rightarrow$  *bounds*”

**fun** *plus\_bounds* :: “*bounds*  $\Rightarrow$  *bounds*  $\Rightarrow$  *bounds*”

**global\_interpretation** *Val\_semilattice*

**where**  $\gamma = \gamma\_bounds$  **and**  $num' = num\_bounds$  **and**  $plus' = plus\_bounds$

**global\_interpretation** *Abs\_Int*

**where**  $\gamma = \gamma\_bounds$  **and**  $num' = num\_bounds$  **and**  $plus' = plus\_bounds$

**defines**  $aval\_bounds = aval'$  **and**  $step\_bounds = step'$  **and**  $AI\_bounds = AI$

**global\_interpretation** *Abs\_Int\_mono*

**where**  $\gamma = \gamma\_bounds$  **and**  $num' = num\_bounds$  **and**  $plus' = plus\_bounds$

**fun** *m\_bounds* :: “*bounds*  $\Rightarrow$  *nat*”

**abbreviation** *h\_bounds* :: *nat*

**global\_interpretation** *Abs\_Int\_measure*

**where**  $\gamma = \gamma\_bounds$  **and**  $num' = num\_bounds$  **and**  $plus' = plus\_bounds$

**and**  $m = m\_bounds$  **and**  $h = h\_bounds$